
Instructions: Language of the Computer

Computer Architecture

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Six Steps in Execution of a Procedure

1. Main routine(**caller**) places parameters in a place where the procedure (**callee**) can access them
 - \$a0 - \$a3 : four **argument** registers
2. **Caller** transfers control to the **callee**
3. **Callee** acquires the storage resources needed
4. **Callee** performs the desired task
5. **Callee** places the result value in a place where the **caller** can access it
 - \$v0 - \$v1 : two **value** registers for result values
6. **Callee** returns control to the **caller**
 - \$ra : one **return address** register to return to the point of origin

Register Usage

- \$a0 – \$a3: arguments (reg's 4 – 7)
- \$v0, \$v1: result values (reg's 2 and 3)
- \$t0 – \$t9: temporaries
 - Can be overwritten by callee
- \$s0 – \$s7: saved
 - Must be saved/restored by callee
- \$gp: global pointer for static data (reg 28)
- \$sp: stack pointer (reg 29)
- \$fp: frame pointer (reg 30)
- \$ra: return address (reg 31)

Procedure Call Instructions

- Procedure call: jump and link
`jal ProcedureLabel`
 - Address of following instruction put in `$ra`
 - Jumps to target address
- Procedure return: jump register
`jr $ra`
 - Copies `$ra` to program counter
 - Can also be used for computed jumps
 - e.g., for case/switch statements

Leaf Procedure Example

- C code:

```
int leaf_example (int g, h, i, j)
{ int f;
  f = (g + h) - (i + j);
  return f;
}
```

- Arguments g, ..., j in \$a0, ..., \$a3
- f in \$s0 (hence, need to save \$s0 on stack)
- Result in \$v0

Leaf Procedure Example

- MIPS code:

leaf_example:			
addi	\$sp,	\$sp,	-4
sw	\$s0,	0(\$sp)	
add	\$t0,	\$a0,	\$a1
add	\$t1,	\$a2,	\$a3
sub	\$s0,	\$t0,	\$t1
add	\$v0,	\$s0,	\$zero
lw	\$s0,	0(\$sp)	
addi	\$sp,	\$sp,	4
jr	\$ra		

Save \$s0 on stack

Procedure body

Result

Restore \$s0

Return

Non-Leaf Procedures

- Procedures that call other procedures
- For nested call, caller needs to save on the stack:
 - Its return address
 - Any arguments and temporaries needed after the call
- Restore from the stack after the call

Non-Leaf Procedure Example

- C code:

```
int fact (int n)
{
    if (n < 1) return 1;
    else return n * fact(n - 1);
}
```

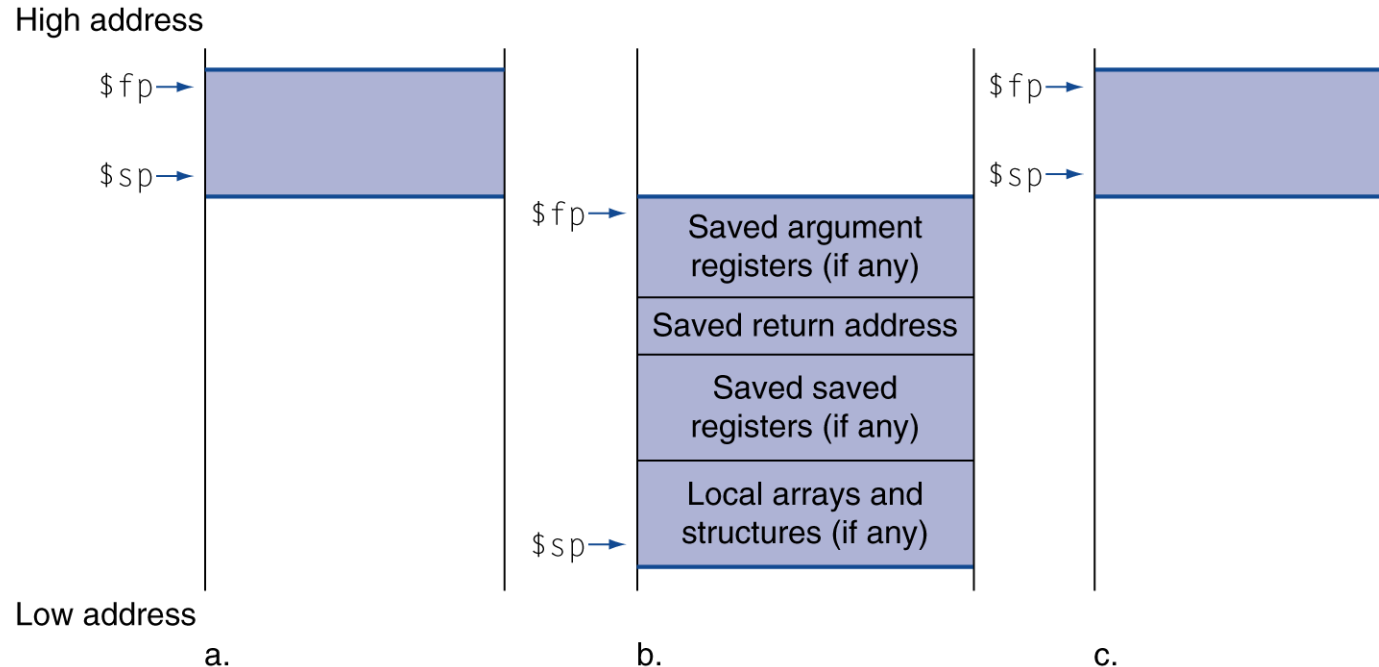
- Argument n in \$a0
- Result in \$v0

Non-Leaf Procedure Example

- MIPS code:

fact:		
	addi \$sp, \$sp, -8	# adjust stack for 2 items
	sw \$ra, 4(\$sp)	# save return address
	sw \$a0, 0(\$sp)	# save argument
	slti \$t0, \$a0, 1	# test for n < 1
	beq \$t0, \$zero, L1	
	addi \$v0, \$zero, 1	# if so, result is 1
	addi \$sp, \$sp, 8	# pop 2 items from stack
	jr \$ra	# and return
L1:	addi \$a0, \$a0, -1	# else decrement n
	jal fact	# recursive call
	lw \$a0, 0(\$sp)	# restore original n
	lw \$ra, 4(\$sp)	# and return address
	addi \$sp, \$sp, 8	# pop 2 items from stack
	mul \$v0, \$a0, \$v0	# multiply to get result
	jr \$ra	# and return

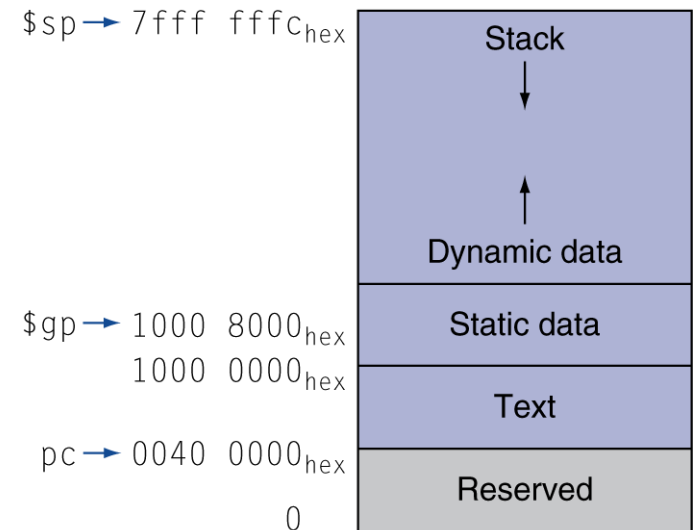
Local Data on the Stack



- Local data allocated by callee
 - e.g., C automatic variables
- Procedure frame (activation record)
 - Used by some compilers to manage stack storage

Memory Layout

- Text: program code
- Static data: global variables
 - e.g., static variables in C, constant arrays and strings
 - \$gp initialized to address allowing \pm offsets into this segment
- Dynamic data: heap
 - E.g., malloc in C, new in Java
- Stack: automatic storage



Character Data

- Byte-encoded character sets
 - ASCII: 128 characters
 - 95 graphic, 33 control
 - Latin-I: 256 characters
 - ASCII, +96 more graphic characters
- Unicode: 32-bit character set
 - Used in Java, C++ wide characters, ...
 - Most of the world's alphabets, plus symbols
 - UTF-8, UTF-16: variable-length encodings

Byte/Halfword Operations

- Could use bitwise operations
- MIPS byte/halfword load/store
 - String processing is a common case

lb rt, offset(rs) lh rt, offset(rs)

- Sign extend to 32 bits in rt

lbu rt, offset(rs) lhu rt,
offset(rs)

- Zero extend to 32 bits in rt

sb rt, offset(rs) sh rt, offset(rs)

- Store just rightmost byte/halfword