
The Processor

Computer Architecture

2019 I 학기

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홍석준

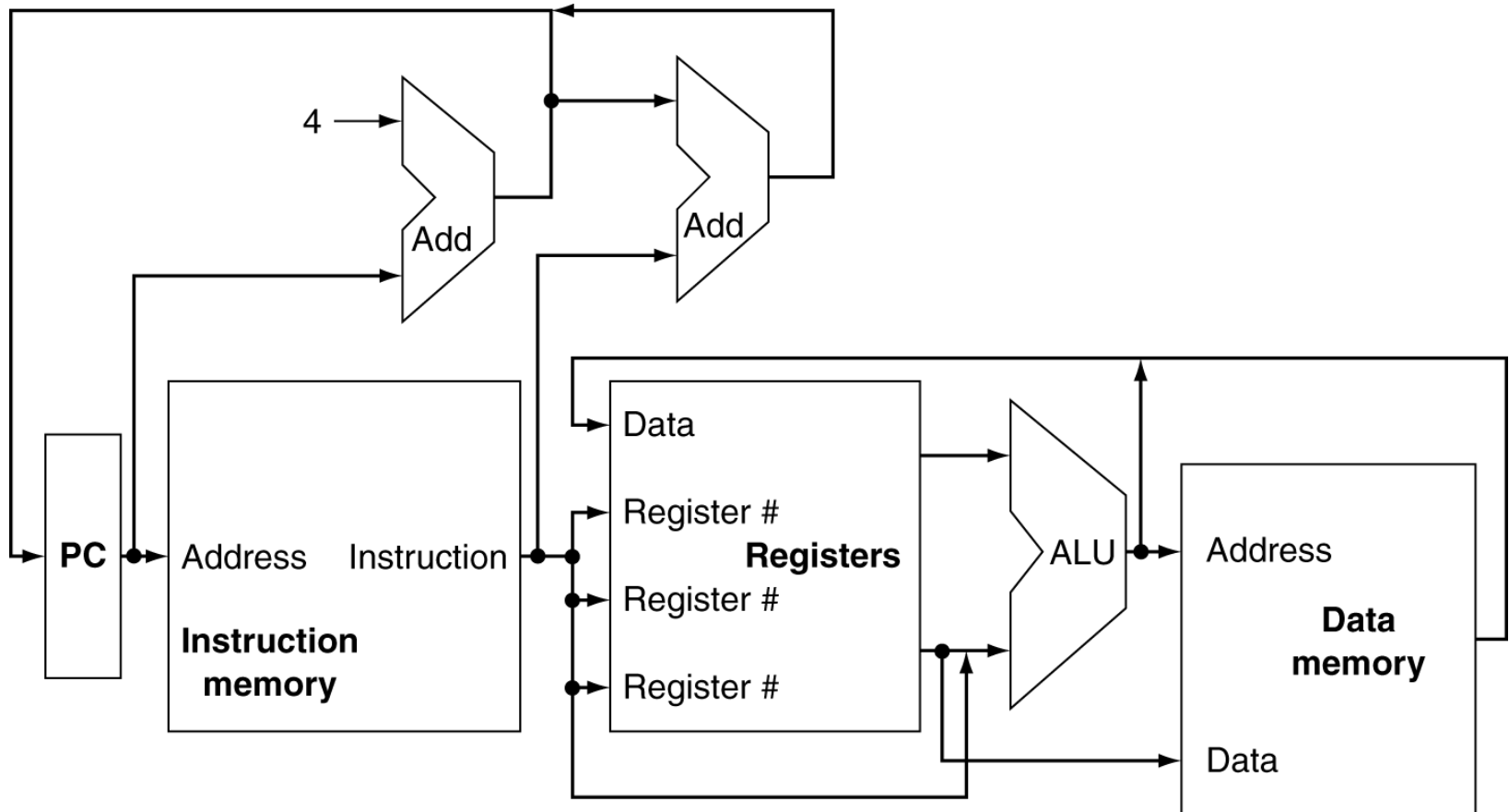
Introduction

- CPU performance factors
 - Instruction count
 - Determined by ISA and compiler
 - CPI and Cycle time
 - Determined by CPU hardware
- We will examine two MIPS implementations
 - A simplified version
 - A more realistic pipelined version
- Simple subset, shows most aspects
 - Memory reference: lw, sw
 - Arithmetic/logical: add, sub, and, or, slt
 - Control transfer: beq, j

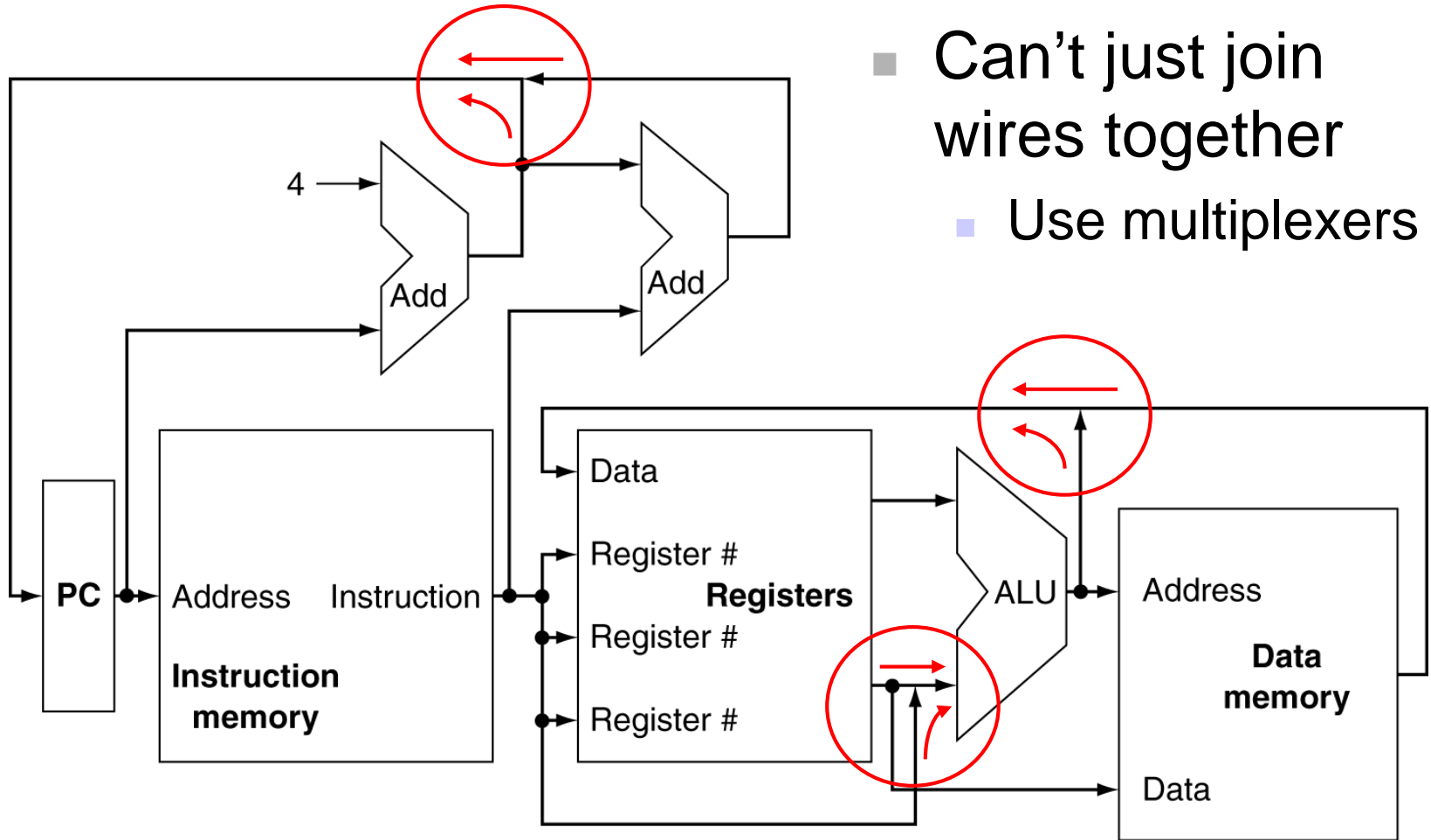
Instruction Execution

- PC → instruction memory, fetch instruction
- Register numbers → register file, read registers
- Depending on instruction class
 - Use ALU to calculate
 - Arithmetic result
 - Memory address for load/store
 - Branch target address
 - Access data memory for load/store
 - $PC \leftarrow \text{target address or } PC + 4$

CPU Overview

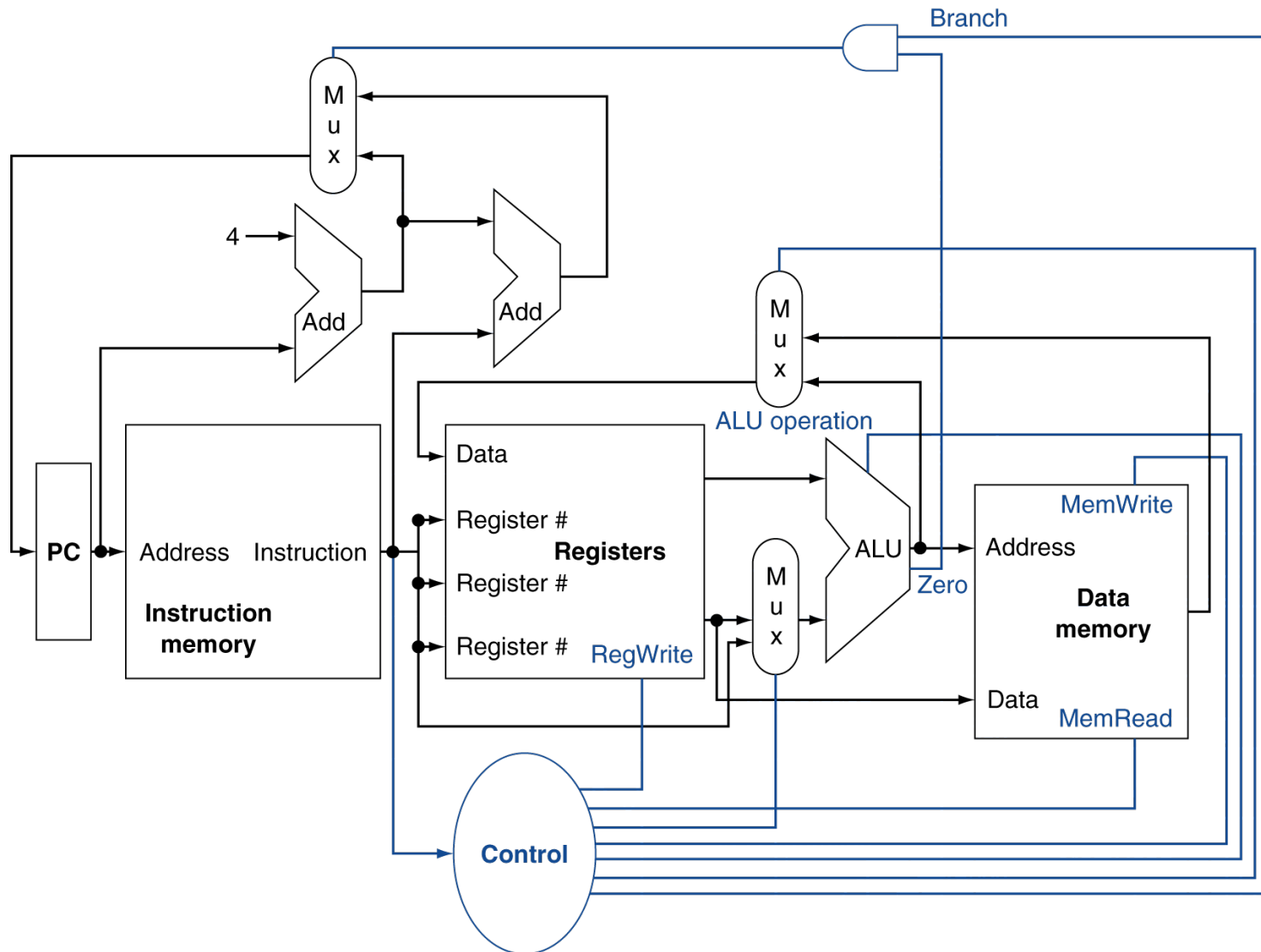


Multiplexers



- Can't just join wires together
 - Use multiplexers

Control



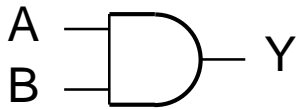
Logic Design Basics

- Information encoded in binary
 - Low voltage = 0, High voltage = 1
 - One wire per bit
 - Multi-bit data encoded on multi-wire buses
- Combinational element
 - Operate on data
 - Output is a function of input
- State (sequential) elements
 - Store information

Combinational Elements

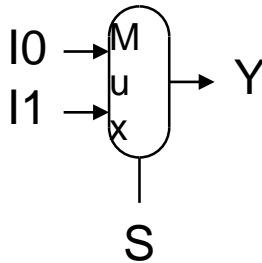
- AND-gate

- $Y = A \& B$



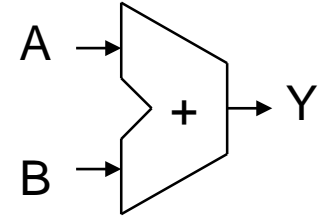
- Multiplexer

- $Y = S ? I1 : I0$



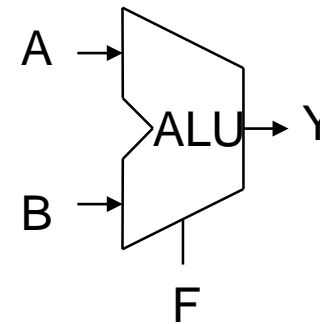
- Adder

- $Y = A + B$



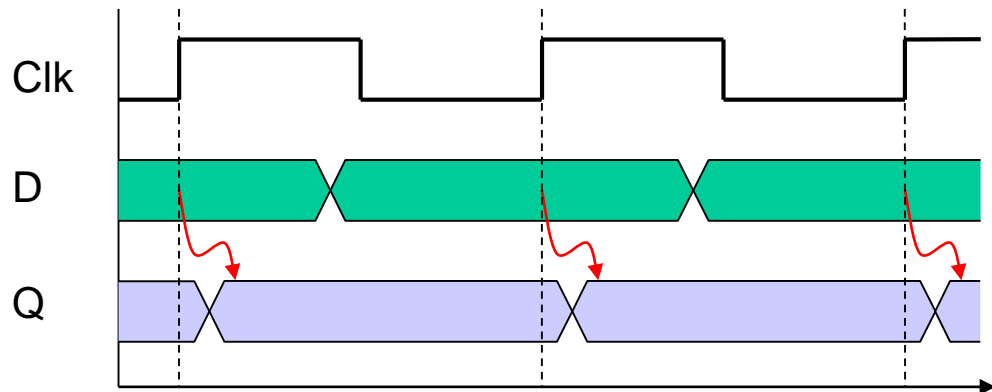
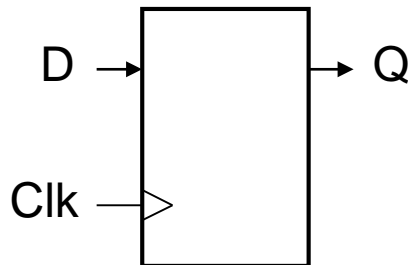
- Arithmetic/Logic Unit

- $Y = F(A, B)$



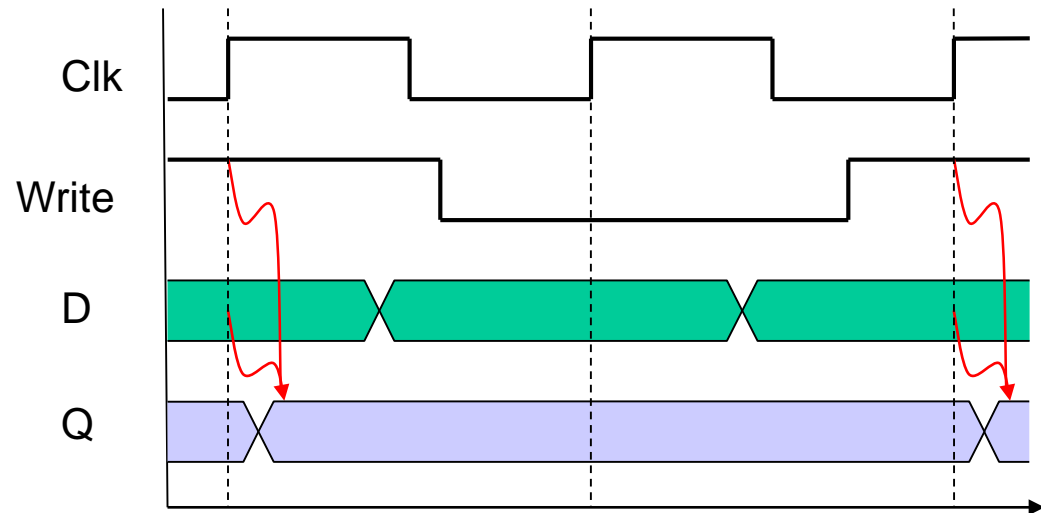
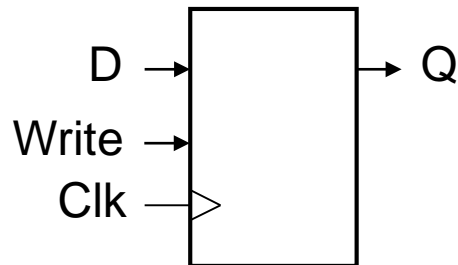
Sequential Elements

- Register: stores data in a circuit
 - Uses a clock signal to determine when to update the stored value
 - Edge-triggered: update when Clk changes from 0 to 1



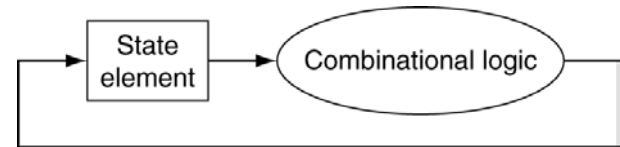
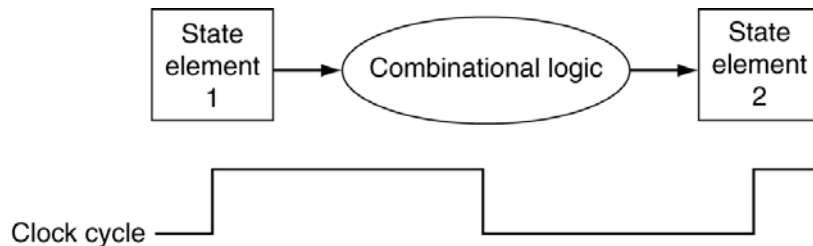
Sequential Elements

- Register with write control
 - Only updates on clock edge when write control input is 1
 - Used when stored value is required later



Clocking Methodology

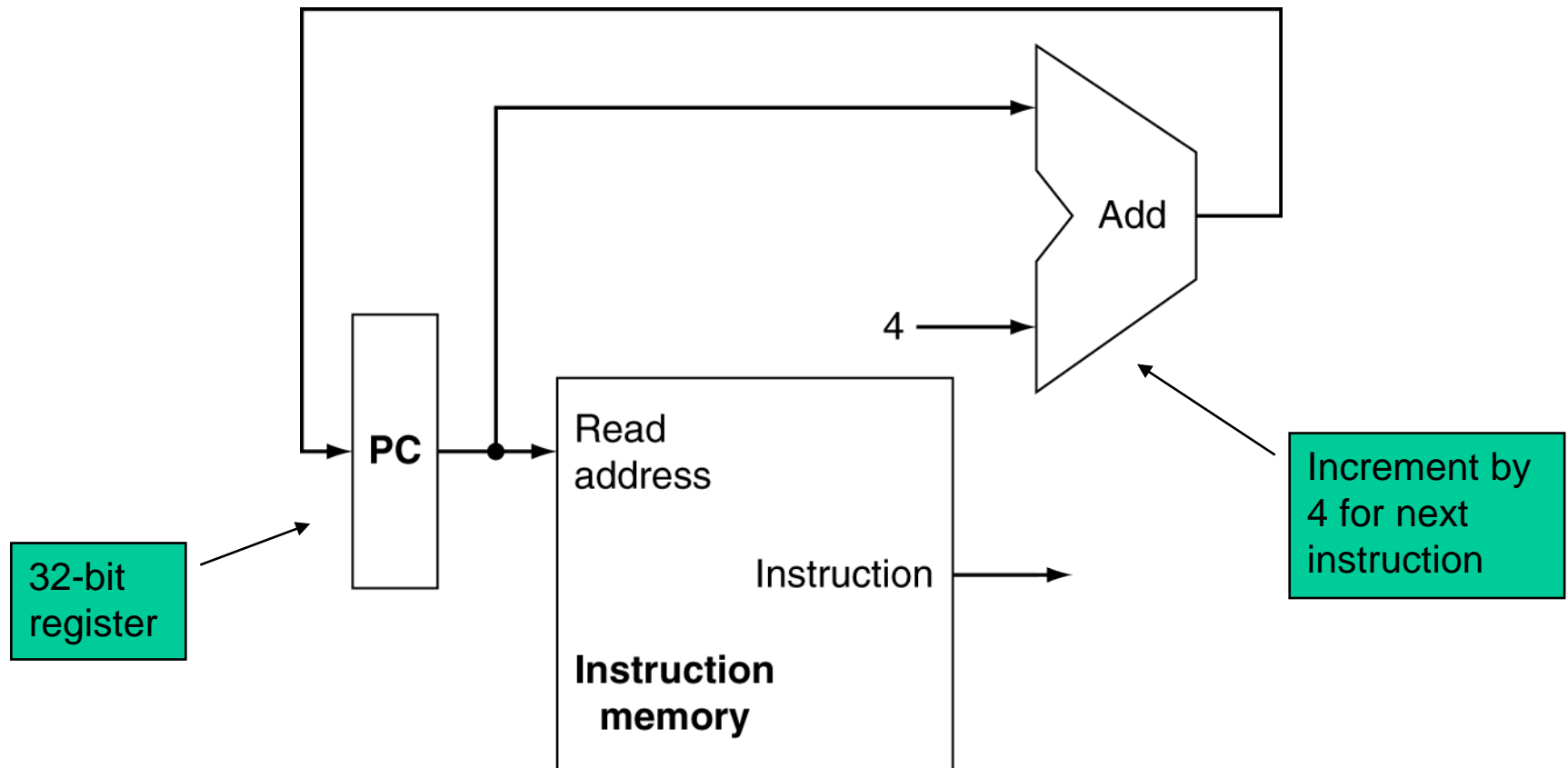
- Combinational logic transforms data during clock cycles
 - Between clock edges
 - Input from state elements, output to state element
 - Longest delay determines clock period



Building a Datapath

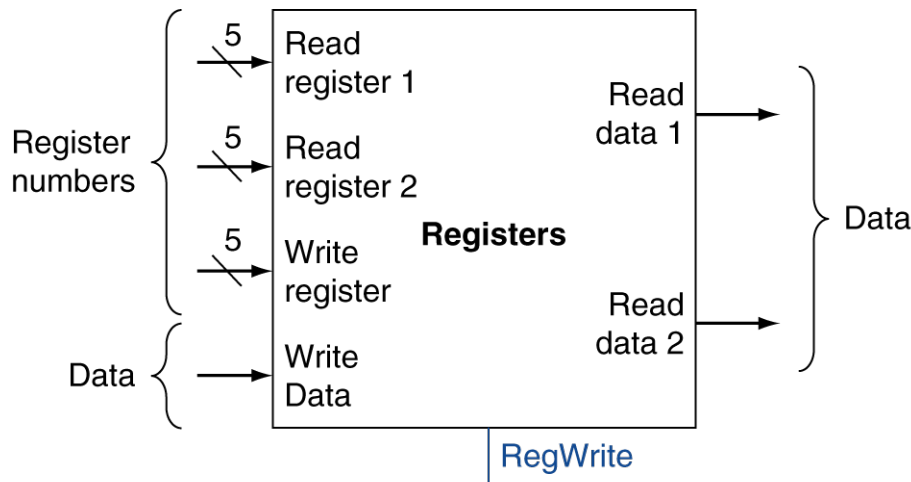
- Datapath
 - Elements that process data and addresses in the CPU
 - Registers, ALUs, mux's, memories, ...
- We will build a MIPS datapath incrementally
 - Refining the overview design

Instruction Fetch

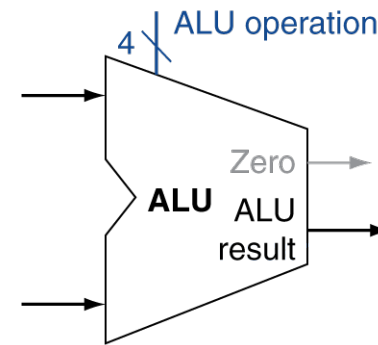


R-Format Instructions

- Read two register operands
- Perform arithmetic/logical operation
- Write register result



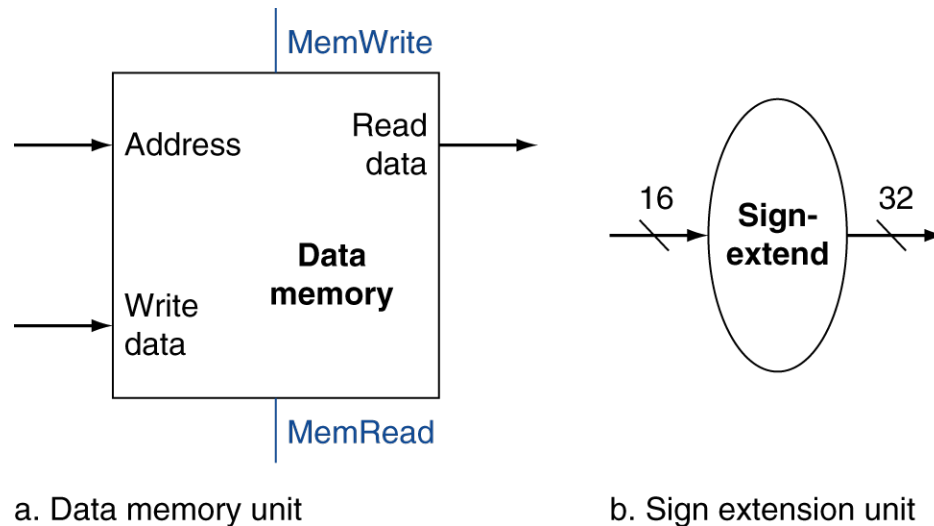
a. Registers



b. ALU

Load/Store Instructions

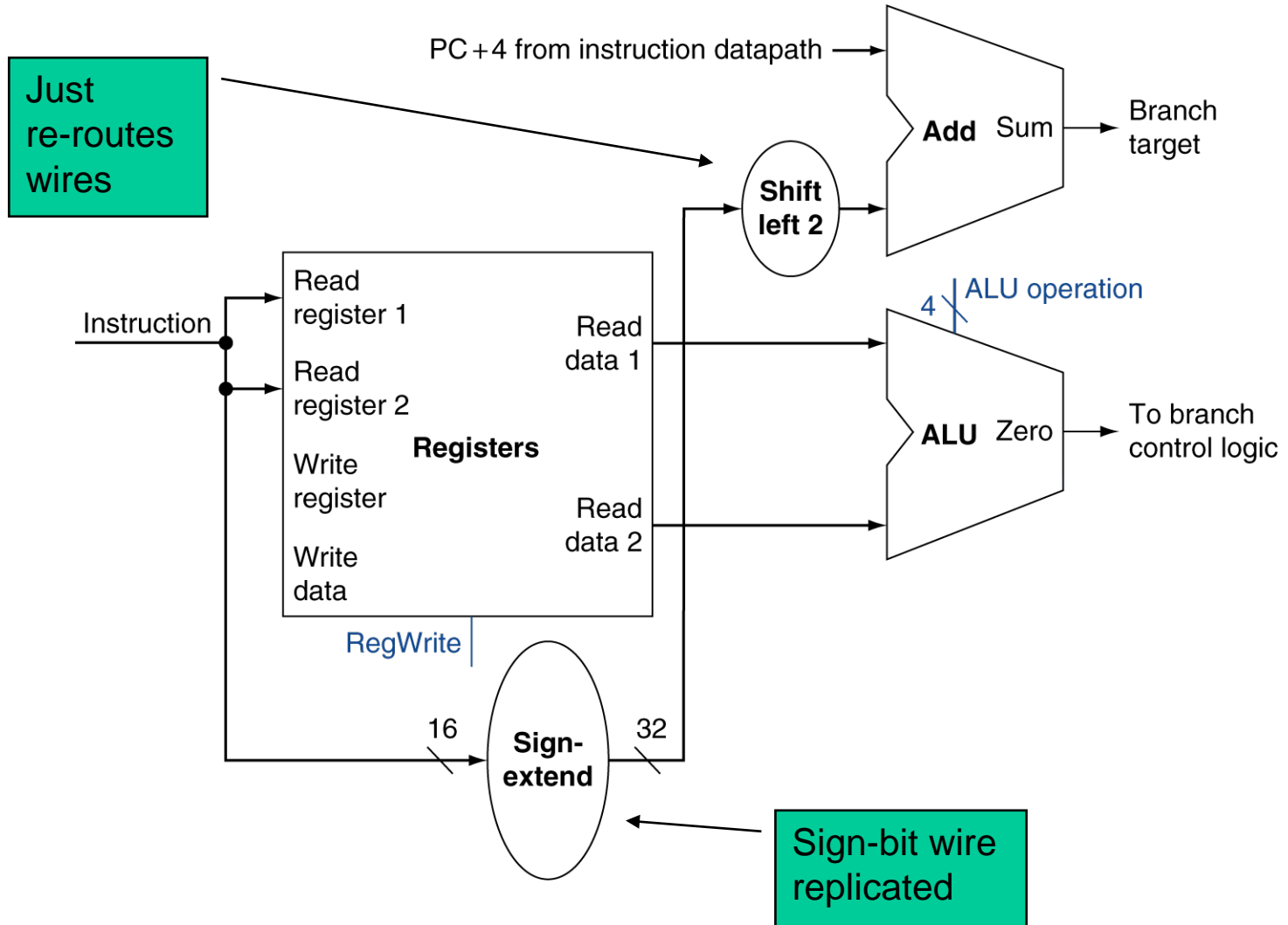
- Read register operands
- Calculate address using 16-bit offset
 - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory



Branch Instructions

- Read register operands
- Compare operands
 - Use ALU, subtract and check Zero output
- Calculate target address
 - Sign-extend displacement
 - Shift left 2 places (word displacement)
 - Add to PC + 4
 - Already calculated by instruction fetch

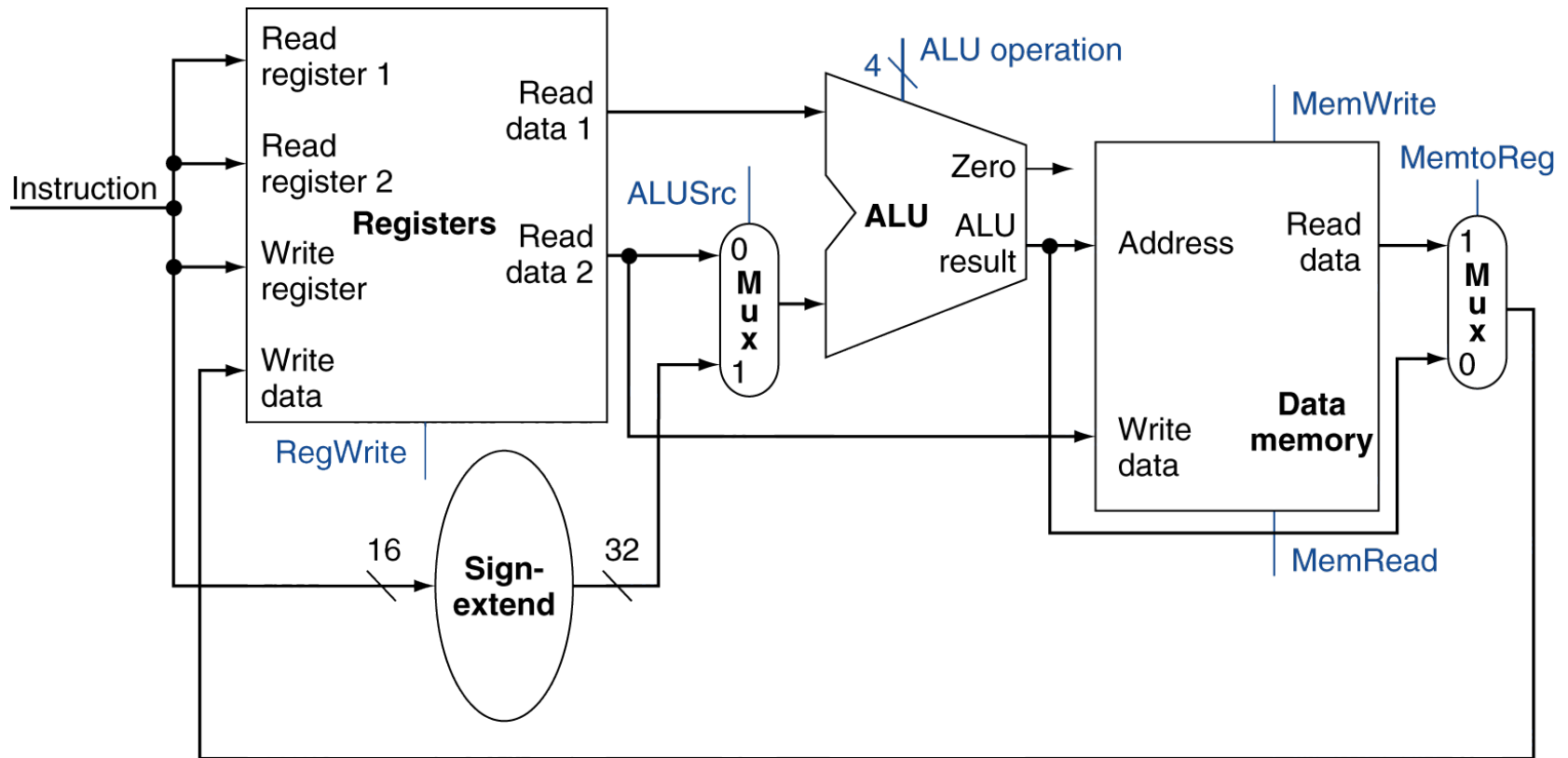
Branch Instructions



Composing the Elements

- First-cut data path does an instruction in one clock cycle
 - Each datapath element can only do one function at a time
 - Hence, we need separate instruction and data memories
- Use multiplexers where alternate data sources are used for different instructions

R-Type/Load/Store Datapath



Full Datapath

